Parallel Transactions

GeneXus

Customer X				\sim
Structure 🥌 Web Layout Rule	es Events Variables Patterns			
ame	Type	Description	Formula	Nullable
Customer	Customer	Customer		
P CustomerId	Id	Customer Id		No
- P CustomerName	Character(50)	Customer Name		No
CustomerAddress	Address, GeneXus	Customer Address		No
Customer × Technical	×			~
tructure 🧏 Web Layout Rule	es Events Variables Patterns			
ame	Туре	Description	Formula	Nullable
Technical	Technical	Technical		
P TechnicalId	Id	Technical Id		No
P TechnicalName	Character(50)	Technical Name		No
Customer X 📑 Technical	X 📑 Material X			~
web Layour Kuie	s Evenis Variables Patterns	Description	Formula	No. Hobble
anne	Type	Description	ronnula	nuidDie
Material	Material	Material		
To Materialid	10	material 1d		NO

Let's suppose that we are designing a small application for a company that provides technical services at the customers' location. From this application, among other things, we will be able to generate the orders for repair requests made by customers.

Let's look only at the transactions we're interested in for our example. We have a transaction to record customers; another transaction to record the company's technicians; and a transaction where the materials can be recorded, which the technicians use to solve the different issues of clients.

RepairOrder Trn

Structure

Structure Structure Kules E	vents Variables Patterns	
Name	Туре	Description
E RepairOrder	RepairOrder	Repair Order
- P RepairOrderId	Id	Repair Order Id
- > CustomerId	Id	Customer Id
- 🖌 CustomerName	Character(50)	Customer Name
 CustomerAddress 	Address, GeneXus	Customer Address
- P RepairOrderEnteredDate	Date	Repair Order Entered Date
 RepairOrderType 	OrderType	Repair Order Type
 RepairOrderPrice 	Price	Repair Order Price
 RepairOrderStatus 	Status	Repair Order Status
 RepairOrderScheduledDate 	Date	Repair Order Scheduled Date
RepairOrderDescription	LongVarChar(2M)	Repair Order Description



Then we have this transaction (RepairOrder), where the operator who receives the request from the customer will enter all the necessary data to generate the service order.

In it, we have as a primary key the RepairOrderId attribute, which has an autonumbered data type assigned to it. Each order will have an associated client, so in this transaction we have inferred attributes from the Customer table. Lastly, we have these other attributes of the transaction. To record the date of entry, the type of order (whose data type will be enumerated containing these values), the price, a status (whose data type will also be enumerated, and will accept the following values), the date scheduled and a description.

Let's look at the rules we entered.

We have a couple of "default" rules and several "noaccept" rules. In the default rules we assign a value to the attribute that will keep the status of the order –in this case "P", for pending. And in this other case, we will assign today's date to the attribute that will save precisely the date of entry of the work order.

We also see that this attribute (RepairOrderPrice) of the transaction will be calculated from the value entered in this other attribute (RepairOrderType), which will record the type of order.

Lastly, we declare a couple of "error" rules.

Let's see this transaction running, from which we enter a repair order.



Now let's suppose that in our reality, after an order is entered, a coordinator must assign a technician to it and the materials provided to that technician, in order to monitor it. And change its status to assigned.

In addition, when the task is completed by the technician, an administrative assistant should be able to assign an end date to the order, enter comments from the technician and the customer, record whether the issue was solved or not, and change the status of the order to completed.

How can we do this?

RepairOrder Trn with all attributes

	and an and a second the			
me	Туре	Description	Formula	Nullable
RepairOrder	RepairOrder	Repair Order		
- 📍 RepairOrderId	Id	Repair Order Id		No
- 🛃 CustomerId	Id	Customer Id		No
 <i>L</i>ustomerName 	Character(50)	Customer Name		
CustomerAddress	Address, GeneXus	Customer Address		
- P RepairOrderEnteredDate	Date	Repair Order Entered Date		No
 RepairOrderType 	OrderType	Repair Order Type		No
 RepairOrderPrice 	Price	Repair Order Price		No
 RepairOrderStatus 	Status	Repair Order Status		No
 RepairOrderScheduledDate 	Date	Repair Order Scheduled Date		No
 RepairOrderDescription 	LongVarChar(2M)	Repair Order Description		No
- 🛃 TechnicalId	Id	Technical Id		Yes
- 🖌 TechnicalName	Character(50)	Technical Name		
 RepairOrderFinalizedDate 	Date	Repair Order Finalized Date		No
RepairOrderTecObservations	LongVarChar(2M)	Repair Order Tec Observations		No
RepairOrderCustObservations	LongVarChar(2M)	Repair Order Cust Observations		No
RepairOrderProblemSolved	Boolean	Repair Order Problem Solved		No
🗃 🖢 Material	Material	Material		
P MaterialId	Id	Material Id		No
MaterialDescpription	Numeric(4.0)	Material Descpription		No
 MaterialQty 	Numeric(4.0)	Material Oty		No

One option would be to add all these attributes to the transaction, and to complete those that we need at every moment. The transaction would look like this.

Let's see it at runtime.

As we can see, this solution does not provide a friendly interface for the user, since there will always be fields that may not be of interest or may not have to be completed.

For example, when the operator receives the order from the customer and enters it for the first time, before it is assigned to a technician, we will see several fields that will not be completed. In this case they are relatively few, but if the number of attributes is increased this will impact more and more on the users and their experience with the system.

The same happens when the coordinator accesses the transaction, selects the order to be modified and so it is changed to Update mode, assigns a technician to it and loads the materials.

As soon as we search for the order and load it on the screen, we will be shown attributes that we may not be interested in viewing. Then, to load the materials, you will have to leave fields empty in order to complete them later. Of course, this is not comfortable for entering data. And as we've just said, since there are few attributes in this example, it doesn't have a great impact. However, as the number of attributes increases, since we may need to record more data, the input process becomes more complex. When the administrative assistant wants to complete it with the corresponding observations, he will need to select the order to work on, switching the transaction to Update mode, and will see all the attributes of this transaction, when in fact he only needs to see a few details.

Another disadvantage is how complex it can be to program the behavior of this transaction. As we saw, it will go through different states and modifications with later entries. Also, it may require different controls every time, even of the attributes themselves. But, given this design, everything will be programmed in the same object. It will not be possible to customize rules or events specific to each of the moments we have just seen.

These disadvantages could be solved in a simple way by using what we call parallel transactions.

Parallel transactions

RepairOrder X					V RepairOrder X
Structure SWeb Layout Rules Events	Variables Patterns				Structure SWeb Layout Rules Events Variables Patterns
Name	Туре	Description	Formula	Nullable	1 default(RepairOrderStatus, Status.Pending);
Em RepairOrder	RepairOrder	Repair Order			<pre>2 default(RepairOrderEnteredDate, &Today);</pre>
- P RepairOrderId	Id	Repair Order Id		No	3
- > CustomerId	Id	Customer Id		No	4 noaccept(RepairOrderId);
- 🖌 CustomerName	Character(50)	Customer Name			5 noaccept(RepairOrderPrice);
─	Address, GeneXus	Customer Address			6 □ noaccept(RepairOrderEnteredDate);
- P RepairOrderEnteredDate	Date	Repair Order Entered Date		No	
RepairOrderType	OrderType	Repair Order Type		No	S RepairOrderPrice = 500 if RepairOrderType = OrderType.Basic;
RepairOrderPrice	Price	Repair Order Price		No	³ RepairOrderPrice = 1000 if RepairOrderType = OrderType.Intermediate;
 RepairOrderStatus 	Status	Repair Order Status		No	10 p RepairorderPrice = 1500 if RepairorderType = UnderType.Advanced;
 RepairOrderScheduledDate 	Date	Repair Order Scheduled Date		No	12 Gamman("The scheduling data cannot be explice than today") if PanaioDodesScheduleDate .
 RepairOrderDescription 	LongVarChar(2M)	Repair Order Description		No	13 Jacober ('You must enter a description') if Beneficienter Comy ().
					14
RepairOrderAssigned X					✓ []
itructure 5 Web Layout Rules Events	Variables Patterns				RepairOrderAssigned ×
ame	Туре	Description	Formula	Nullable	Structure & Web Laword Budge Screek Visibility Datame
RepairOrderAssigned	RepairOrderAssigned	Repair Order Assigned			Subcure Web Layout Rules Variables Patterns
- Y RepairOrderId	Id	Repair Order Id		No	1 CustomerId.Visible = false;
- > CustomerId	Id	Customer Id		No	2 RepairOrderStatus = Status.Assigned;
CustomerAddress	Address, GeneXus	Customer Address			
- P RepairOrderScheduledDate	Date	Repair Order Scheduled Date		No	<pre>4 Derror('You must enter a technician') if TechnicalId.IsEmpty(); </pre>
 RepairOrderStatus 	Status	Repair Order Status		No	
TechnicalId	Id	Technical Id		Yes	
 ZechnicalName 	Character(50)	Technical Name			<pre>% hoaccept(ustomeria); % moaccept(ustomeria);</pre>
Material	Material	Material			9 maccept (Renain Order Scatch) (athata);
- 📍 MaterialId	Id	Material Id		No	10
 MaterialDescription 	LongVarChar(2M)	Material Description			
MaterialQty	Numeric(4.0)	Material Qty		No	1.
RepairOrderCompleted X					V RepairOrderCompleted X
Structure 🐔 Web Layout Rules Events	Variables Patterns				Structure S Web Layout Rules Events Variables Patterns
ame	Туре	Description	Formula	Nullable	1 RepairOrderStatus = status.Finalized;
RepairOrderCompleted	RepairOrderCompleted	Repair Order Completed			2
- P RepairOrderId	Id	Repair Order Id		No	Incaccept(ReparOrderStatus);
- P RepairOrderFinalizedDate	Date	Repair Order Finalized Date		No	<pre>%pidetauit(RepairOrder>inalizedUate, &Today); %</pre>
 RepairOrderStatus 	Status	Repair Order Status		No	formed 'Did ant enter technician observations') if Pennished. The
 RepairOrderTecObservations 	LongVarChar(2M)	Repair Order Tec Observations		No	7 msg(big not enter technician Observations') if RepairOrderTecObservations.IstBpty(); 7 msg(big entermain entermatic
RepairOrderCustObservations	LongVarChar(2M)	Repair Order Cust Observations		No	<pre>*pinst no customer comments encered) if RepairOrderCustosservations.isEmpty(); 8</pre>

In this way, we would have a transaction for entering the order first by the operator. Another transaction, where the coordinator will assign that order to the technician and load the materials to be given to him. Lastly, a transaction, where the administrative assistant will be able to upload the technicians' and customer's comments, if any, and whether the problem was solved or not. Each one of these transactions will have the attributes that we are really interested in for each case.

For this, all three transactions must have the same attribute as the primary key. This is what makes them parallel transactions.

It is important to mention that the parallelization of transactions is done by level. In this example, the main levels of these transactions are parallel.

If we look at the rules section of the transactions, we see that each one has its own rules, totally independent from each other. This is one of the great advantages of working with parallel transactions. The same is true for any events we may program.

Let's see the application running with these modifications.

	Browse Web Objec	ts
Customer RepairOrder Techacal	Material	RepairOrder
OPERATOR	COORDINATOR	ADMINISTRATIVE ASSISTANT

From this transaction (RepairOrder) an operator receives the customer's request and completes all the fields. We see that the value "P" –of pending– is assigned automatically, because it is defined in the default rule.

Some time later, the coordinator accesses this other transaction (RepairOrderAssigned), searches for the order, and completes the necessary data. The value that records the status of the order is automatically changed to "A" (assigned).

Finally, after the task is completed and the technical report is received, the administrator will manage it from the third transaction created (RepairOrderCompleted), filling in the corresponding fields, and the status will be automatically changed assigning it the value "F" (finished).

We see that, in this way, it becomes much simpler, clearer and more intuitive for users. And it helps to make the application scalable, to face future changes that may have to be made in the procedure for entering orders.

In Database



At the database level, from the main levels of these three transactions, a single physical table will be generated, since as we know, GeneXus designs the database following normalization criteria.

The generated table will contain the attributes resulting from associating the three transactions.

When we enter a record from the first transaction, what happens then in the table of our database, with the attributes to which we haven't assigned any value yet? What value will they have?

Empty values will be assigned by default: "0" in numeric fields, empty string in character type fields, a default date representing the empty value for those of Date type, and false for Boolean attributes.

Parallel transactions offer a clear way to keep in the structure of each transaction solely the information you want to work with within that program, regardless of the information in the same table that can be handled in another transaction.

There is a wide range of realities that can lead us to use this type of transaction, which we will not examine in detail in this video. For more details on this topic, please visit our Wiki.



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